

LINEAR PRICING MECHANISMS FOR MARKETS WITHOUT CONVEXITY

Paul Milgrom¹ and Mitchell Watt²

Setting

Two-sided, quasilinear markets with L commodities
 N buyers, $F = \phi N$ sellers, with fixed costs or other non-convexities

Challenges

1. Computational complexity
2. Balancing budgets
3. Achieving individual rationality
4. Ensuring truthful reporting despite rationing

Walrasian equilibria
 may not exist!

Two new extensions of the Walrasian mechanism

Markup Mechanism (α, p, ω)

Sellers paid prices p and buyers pay prices $(1 + \alpha)p$, allocation ω in supply and demand sets at respective prices, while enforcing physical and budget feasibility

for markets
 with two-sided
 non-convexities

Rationing Mechanism (p, ω)

A single price p is used for payments, but some agents are **rationed** to ensure ω has no unallocated supply

for markets
 with one
 convex side

Bound-form First Welfare Theorem

Given **any** feasible allocation $\omega = (x, y)$, **any** nonnegative prices $p \in \mathbb{R}_+^L$, the surplus function $S: \Omega \rightarrow \mathbb{R}$, and an efficient $\omega^* \in \Omega$,

$$S(\omega^*) - S(\omega) \leq p \cdot \left(\sum_f y_f - \sum_n x_n \right) + \left(\sum_f \mathcal{R}_f(p, y_f) + \sum_n \mathcal{R}_n(p, x_n) \right)$$

loss at $\omega \leq$ budget deficit + sum of **rationing losses**

An agent's **rationing loss** at price p and allocation ω is the difference between their *maximal* payoff at p and their *realized* payoff at (p, ω) .

If Walrasian equilibria exist, both coincide with Walrasian mechanism.

We focus on two simple-to-compute instances of these mechanisms requiring only *convex* optimizations and a single binary search.

Efficiency

Markup mechanism: can choose $\alpha \sim O(1/N)$ leading to $O(1/N)$ percentage welfare loss

Rationing mechanism: can identify p with $O(1/N)$ total loss under strong monotonicity assumption

Incentives

Both mechanisms have similar incentive properties to Walrasian mech, but rationing mech may be $\varepsilon - IR$